

Backing Backtracking

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22nd International Conference on Theory and Applications of Satisfiability Testing (SAT 2019)

July 12, 2019



CDCL Invariants

- Trail:** The assignment trail contains neither complementary pairs of literals nor duplicates.
- ConflictLower:** The assignment trail preceding the current decision level does not falsify the formula.
- Propagation:** On every decision level preceding the current decision level all unit clauses are propagated until completion.
- LevelOrder:** The literals are ordered on the assignment trail in ascending order with respect to their decision level.
- ConflictingClause:** At decision levels greater than zero the conflicting clause contains at least two literals with the current decision level.

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Combining CDCL with Chronological Backtracking

τ	...	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19
l	...	4	5	30	47	15	18	6	-7	-8	45	9	38	-23	17	44	-16
δ	...	3	4	4	4	4	4	5	5	5	5	6	6	6	6	6	6

Combining CDCL with Chronological Backtracking

τ	...	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19
l	...	4	5	30	47	15	18	6	-7	-8	45	9	38	-23	17	44	-16
δ	...	3	4	4	4	4	4	5	5	5	5	6	6	6	6	6	6

decision literal

Combining CDCL with Chronological Backtracking

τ	...	4		5	6	7	8	9		10	11	12	13		14	15	16	17	18	19
l	...	4		5	30	47	15	18		6	-7	-8	45		9	38	-23	17	44	-16
δ	...	3		4	4	4	4	4		5	5	5	5		6	6	6	6	6	6

`block(l , 4)`

Combining CDCL with Chronological Backtracking

τ	...	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19
l	...	4	5	30	47	15	18	6	-7	-8	45	9	38	-23	17	44	-16
δ	...	3	4	4	4	4	4	5	5	5	5	6	6	6	6	6	6

$\text{slice}(l, 4)$

Combining CDCL with Chronological Backtracking

τ	...	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19
l	...	4	5	30	47	15	18	6	-7	-8	45	9	38	-23	17	44	-16
δ	...	3	4	4	4	4	4	5	5	5	5	6	6	6	6	6	6

$$l \leq 4$$

Combining CDCL with Chronological Backtracking

τ	...	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19
l	...	4	5	30	47	15	18	6	-7	-8	45	9	38	-23	17	44	-16
δ	...	3	4	4	4	4	4	5	5	5	5	6	6	6	6	6	6



Combining CDCL with Chronological Backtracking

τ	...	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19
l	...	4	5	30	47	15	18	6	-7	-8	45	9	38	-23	17	44	-16
δ	...	3	4	4	4	4	4	5	5	5	5	6	6	6	6	6	6



conflicting { -47, -17, -44 }

conflict level 6

Combining CDCL with Chronological Backtracking

τ	...	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19
l	...	4	5	30	47	15	18	6	-7	-8	45	9	38	-23	17	44	-16
δ	...	3	4	4	4	4	4	5	5	5	5	6	6	6	6	6	6



conflicting { -47, -17, -44 }

learned { -30, -47, -18, 23 }

jump level 4

Combining CDCL with Chronological Backtracking

τ	...	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19
l	...	4	5	30	47	15	18	6	-7	-8	45	9	38	-23	17	44	-16
δ	...	3	4	4	4	4	4	5	5	5	5	6	6	6	6	6	6



conflicting { -47, -17, -44 }

learned { -30, -47, -18, 23 }

backtrack level 5

Combining CDCL with Chronological Backtracking

τ	\dots	4	5	6	7	8	9	10	11	12	13
l	\dots	4	5	30	47	15	18	6	-7	-8	45
δ	\dots	3	4	4	4	4	4	5	5	5	5

Combining CDCL with Chronological Backtracking

τ	...	4	5	6	7	8	9	10	11	12	13
l	...	4	5	30	47	15	18	6	-7	-8	45
δ	...	3	4	4	4	4	4	5	5	5	5

τ	...	4	5	6	7	8	9	10	11	12	13	14
l	...	4	5	30	47	15	18	6	-7	-8	45	23
δ	...	3	4	4	4	4	4	5	5	5	5	4

out of order

Combining CDCL with Chronological Backtracking

τ	...	4	5	6	7	8	9	10	11	12	13
l	...	4	5	30	47	15	18	6	-7	-8	45
δ	...	3	4	4	4	4	4	5	5	5	5

τ	...	4	5	6	7	8	9	10	11	12	13	14
l	...	4	5	30	47	15	18	6	-7	-8	45	23
δ	...	3	4	4	4	4	4	5	5	5	5	4

$\text{block}(l, 4)$

Combining CDCL with Chronological Backtracking

τ	...	4	5	6	7	8	9	10	11	12	13
l	...	4	5	30	47	15	18	6	-7	-8	45
δ	...	3	4	4	4	4	4	5	5	5	5

τ	...	4	5	6	7	8	9	10	11	12	13	14
l	...	4	5	30	47	15	18	6	-7	-8	45	23
δ	...	3	4	4	4	4	4	5	5	5	5	4

$\text{slice}(l, 4)$

Combining CDCL with Chronological Backtracking

τ	...	4	5	6	7	8	9	10	11	12	13
l	...	4	5	30	47	15	18	6	-7	-8	45
δ	...	3	4	4	4	4	4	5	5	5	5

τ	...	4	5	6	7	8	9	10	11	12	13	14
l	...	4	5	30	47	15	18	6	-7	-8	45	23
δ	...	3	4	4	4	4	4	5	5	5	5	4

$$l \leq 4$$

Combining CDCL with Chronological Backtracking

τ	...	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
l	...	18	6	-7	-8	45	23	-38	-9	10	-11	13	16	-17	-25	42	12	-41
δ	...	4	5	5	5	5	4	4	5	6	7	5	4	4	4	4	5	5

Combining CDCL with Chronological Backtracking

τ	...	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
l	...	18	6	-7	-8	45	23	-38	-9	10	-11	13	16	-17	-25	42	12	-41
δ	...	4	5	5	5	5	4	4	5	6	7	5	4	4	4	4	5	5



conflicting { 17, -42, -12 }

Combining CDCL with Chronological Backtracking

τ	...	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
l	...	18	6	-7	-8	45	23	-38	-9	10	-11	13	16	-17	-25	42	12	-41
δ	...	4	5	5	5	5	4	4	5	6	7	5	4	4	4	4	5	5



conflicting { 17, -42, -12 }

backtrack level 4

Combining CDCL with Chronological Backtracking

τ	\dots	4	5	6	7	8	9	10	11
l	\dots	18	23	-38	16	-17	-25	42	-12
δ	\dots	4	4	4	4	4	4	4	4

Calculus

True: $(F, I, \delta) \rightsquigarrow_{\text{True}} \text{SAT}$ if $F|_I = \top$

False: $(F, I, \delta) \rightsquigarrow_{\text{False}} \text{UNSAT}$ if exists $C \in F$ with $C|_I = \perp$ and $\delta(C) = 0$

Unit: $(F, I, \delta) \rightsquigarrow_{\text{Unit}} (F, I\ell, \delta[\ell \mapsto a])$ if $F|_I \neq \top$ and $\perp \notin F|_I$ and exists $C \in F$ with $\{\ell\} = C|_I$ and $a = \delta(C \setminus \{\ell\})$

Jump: $(F, I, \delta) \rightsquigarrow_{\text{Jump}} (F \wedge D, PK\ell, \delta[L \mapsto \infty][\ell \mapsto j])$ if exists $C \in F$ with $PQ = I$ and $C|_I = \perp$ such that $c = \delta(C) = \delta(D) > 0$ and $\ell \in D$ and $\ell|_Q = \perp$ and $F \models D$ and $j = \delta(D \setminus \{\ell\})$ and $b = \delta(P)$ and $j \leq b < c$ and $K = Q_{\leq b}$ and $L = Q_{> b}$

Decide: $(F, I, \delta) \rightsquigarrow_{\text{Decide}} (F, I\ell, \delta[\ell \mapsto d])$ if $F|_I \neq \top$ and $\perp \notin F|_I$ and $\text{units}(F|_I) = \emptyset$ and $V(\ell) \in V$ and $\delta(\ell) = \infty$ and $d = \delta(I) + 1$

Calculus

True: $(F, I, \delta) \rightsquigarrow_{\text{True}} \text{SAT}$ if $F|_I = \top$

False: $(F, I, \delta) \rightsquigarrow_{\text{False}} \text{UNSAT}$ if exists $C \in F$ with $C|_I = \perp$ and $\delta(C) = 0$

Unit: $(F, I, \delta) \rightsquigarrow_{\text{Unit}} (F, I\ell, \delta[\ell \mapsto a])$ if $F|_I \neq \top$ and $\perp \notin F|_I$ and exists $C \in F$ with $\{\ell\} = C|_I$ and $a = \delta(C \setminus \{\ell\})$

Jump: $(F, I, \delta) \rightsquigarrow_{\text{Jump}} (F \wedge D, PK\ell, \delta[L \mapsto \infty][\ell \mapsto j])$ if exists $C \in F$ with $PQ = I$ and $C|_I = \perp$ such that $c = \delta(C) = \delta(D) > 0$ and $\ell \in D$ and $\ell|_Q = \perp$ and $F \models D$ and $j = \delta(D \setminus \{\ell\})$ and $b = \delta(P)$ and $b = c - 1$ and $K = Q_{\leq b}$ and $L = Q_{> b}$

Decide: $(F, I, \delta) \rightsquigarrow_{\text{Decide}} (F, I\ell, \delta[\ell \mapsto d])$ if $F|_I \neq \top$ and $\perp \notin F|_I$ and $\text{units}(F|_I) = \emptyset$ and $V(\ell) \in V$ and $\delta(\ell) = \infty$ and $d = \delta(I) + 1$

Calculus

True: $(F, I, \delta) \rightsquigarrow_{\text{True}} \text{SAT}$ if $F|_I = \top$

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Unit: $(F, I, \delta) \rightsquigarrow_{\text{Unit}} (F, I\ell, \delta[\ell \mapsto a])$ if $F|_I \neq \top$ and $\perp \notin F|_I$ and exists $C \in F$ with $\{\ell\} = C|_I$ and $a = \delta(C \setminus \{\ell\})$

Jump: $(F, I, \delta) \rightsquigarrow_{\text{Jump}} (F \wedge D, PK\ell, \delta[L \mapsto \infty][\ell \mapsto j])$ if exists $C \in F$ with $PQ = I$ and $C|_I = \perp$ such that $c = \delta(C) = \delta(D) > 0$ and $\ell \in D$ and $\ell|_Q = \perp$ and $F \models D$ and $j = \delta(D \setminus \{\ell\})$ and $b = \delta(P)$ and $b = j$ and $K = Q_{\leq b}$ and $L = Q_{> b}$

Decide: $(F, I, \delta) \rightsquigarrow_{\text{Decide}} (F, I\ell, \delta[\ell \mapsto d])$ if $F|_I \neq \top$ and $\perp \notin F|_I$ and $\text{units}(F|_I) = \emptyset$ and $V(\ell) \in V$ and $\delta(\ell) = \infty$ and $d = \delta(I) + 1$

Invariants

Trail: The assignment trail contains neither complementary pairs of literals nor duplicates.

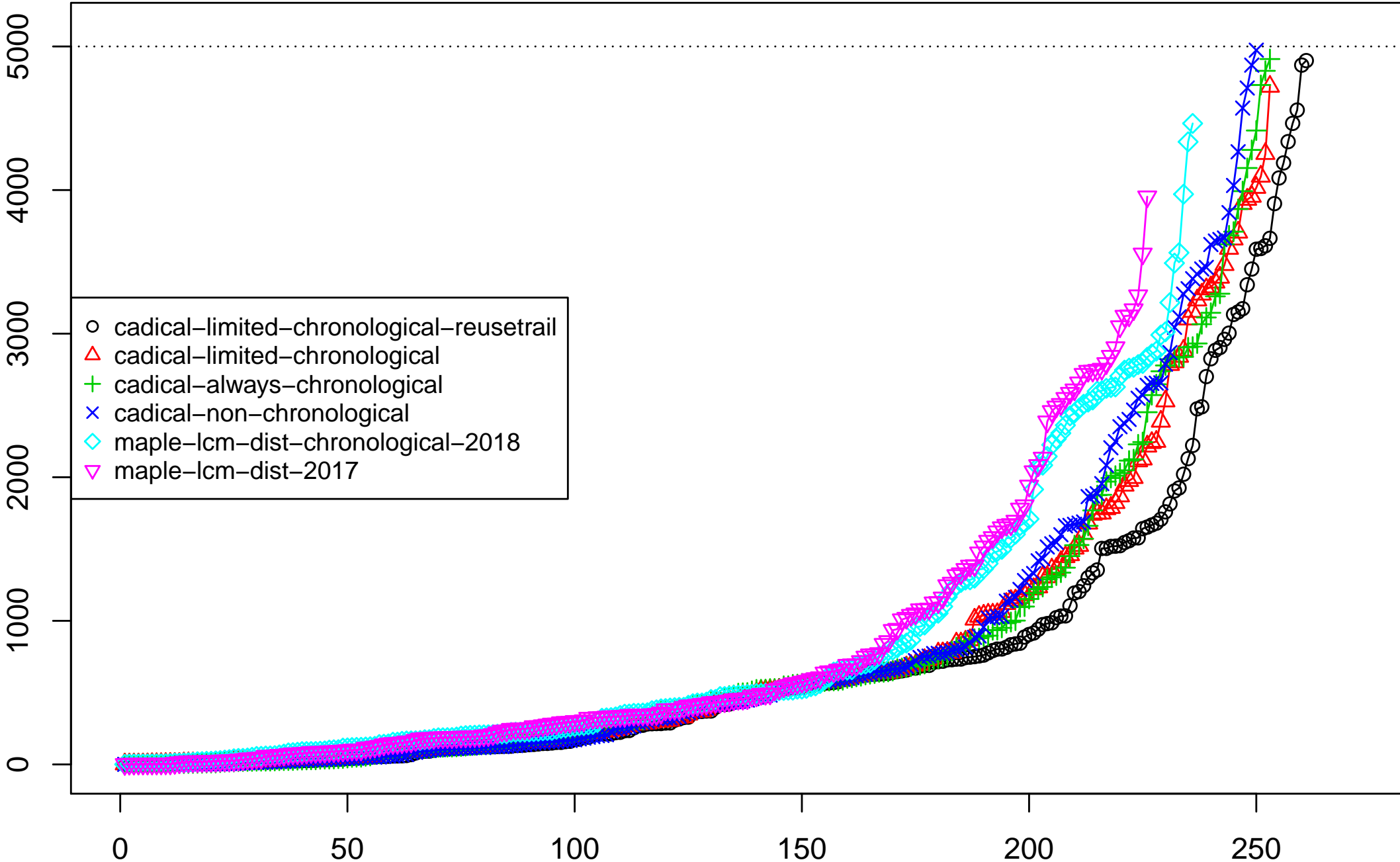
ConflictLower: The assignment trail preceding the current decision level does not falsify the formula.

$$(1): \quad \forall k, \ell \in \text{decs}(I) . \tau(I, k) < \tau(I, \ell) \implies \delta(k) < \delta(\ell)$$

$$(2): \quad \delta(\text{decs}(I)) = \{1, \dots, \delta(I)\}$$

$$(3): \quad \forall n \in \mathbb{N} . F \wedge \text{decs}_{\leq n}(I) \models I_{\leq n}$$

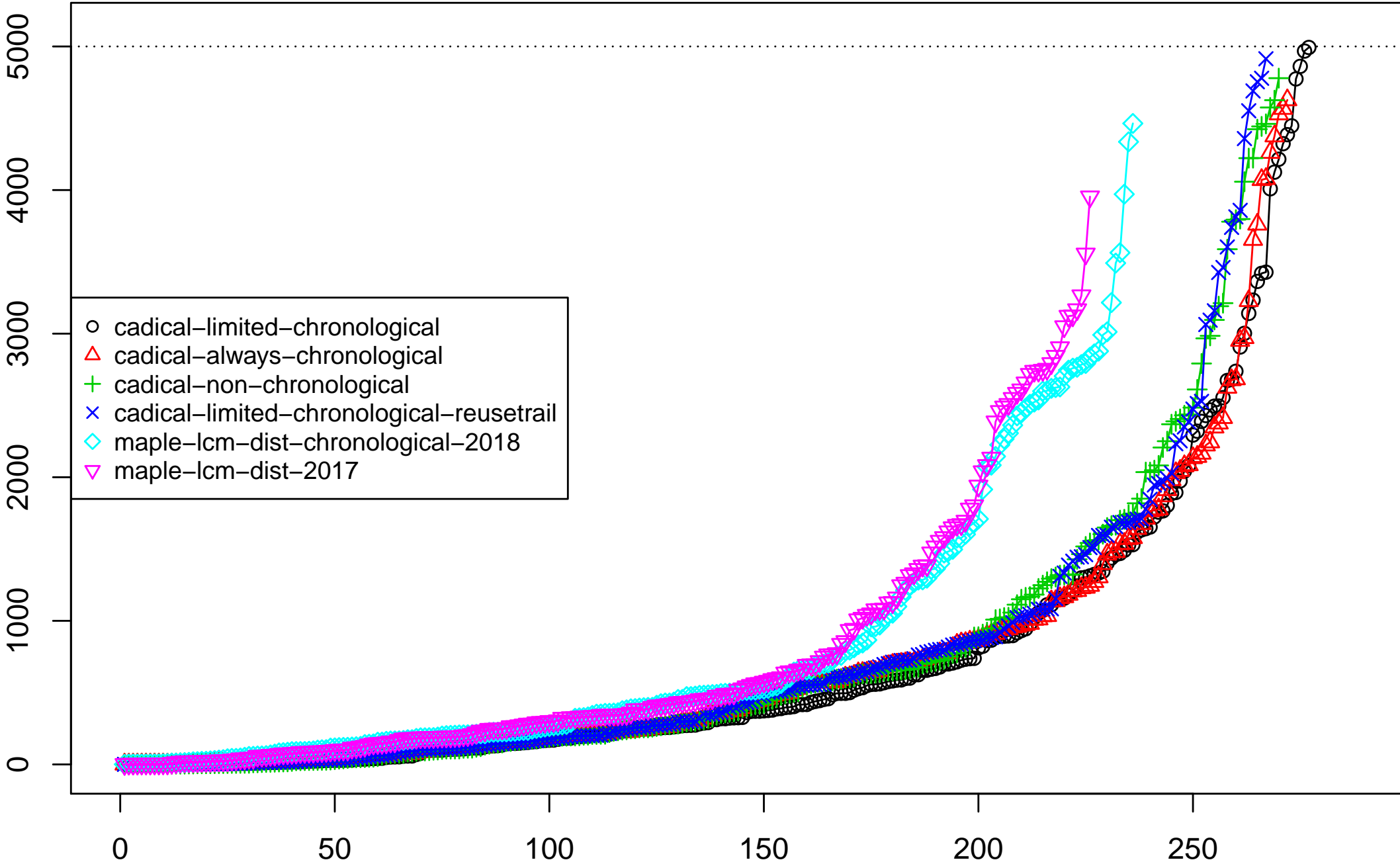
Experiments — Main Track of SAT Competition 2018



Experiments

solver configurations	solved instances		
	total	SAT	UNSAT
cadical-limited-chronological-reusetrail	261	155	106
cadical-limited-chronological	253	147	106
cadical-always-chronological	253	148	105
cadical-non-chronological	250	144	106
maple-lcm-dist-chronological-2018	236	134	102
maple-lcm-dist-2017	226	126	100

Experiments — CaDiCaL Submitted to SAT Race 2019



Experiments – CaDiCaL Submitted to SAT Race 2019

solver configurations	solved instances		
	total	SAT	UNSAT
cadical-limited-chronological	277 (+ 24)	164 (+ 17)	113 (+ 7)
cadical-always-chronological	272 (+ 19)	163 (+ 15)	109 (+ 4)
cadical-non-chronological	270 (+ 20)	160 (+ 16)	110 (+ 4)
cadical-limited-chronological-reusetrail	267 (+ 6)	160 (+ 5)	107 (+ 1)
maple-lcm-dist-chronological-2018	236	134	102
maple-lcm-dist-2017	226	126	100