

## "リ" informatik

## An Abstract CNF-to-d-DNNF Compiler Based on Chronological CDCL

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## The Task: Knowledge Compilation

$$
F \longrightarrow \mathrm{KC}(F) \equiv F
$$

## The Task: CNF-to-d-DNNF Compilation



## Deterministic Decomposable Negation Normal Form

d-DNNF: $\quad F=(a \wedge(b \vee(\neg b \wedge c)) \wedge d) \vee(\neg a \wedge b)$

negations in front of variables
for all conjunctions: conjuncts do not share variables
for all disjunctions: disjuncts are pairwise contradicting

## CNF or d-DNNF - Why Care?



## CNF or d-DNNF - Why Care?



## CNF vs. d-DNNF - the Model Counting Case

CNF: $\quad F=(\neg a \vee b \vee c) \wedge(\neg a \vee d) \wedge(a \vee b)$
d-DNNF: $\quad F^{\prime}=(a \wedge(b \vee(\neg b \wedge c)) \wedge d) \vee(\neg a \wedge b)$

## CNF vs. d-DNNF - the Model Counting Case

CNF: $\quad F=(\neg a \vee b \vee c) \wedge(\neg a \vee d) \wedge(a \vee b)$

$$
\# F=? \quad \text { not that easy }
$$

d-DNNF: $\quad F^{\prime}=(a \wedge(b \vee(\neg b \wedge c)) \wedge d) \vee(\neg a \wedge b)$

$$
\# F^{\prime}=\left(1 \cdot\left(1 \cdot 2^{1}+1 \cdot 1 \cdot 2^{0}\right) \cdot 1\right) \cdot 2^{0}+(1 \cdot 1) \cdot 2^{2}=7
$$

## What About State－of－the－Art Tools？

|  | CDCL | Backtracking | Formula construction | Decomposition |
| :--- | :---: | :---: | :---: | :---: |
| c2d $^{1}$ | $\boldsymbol{v}$ | non－chronological | record trace | static |
| dSharp $^{2}$ | $\boldsymbol{v}$ | non－chronological | record trace | dynamic |
| D4 $^{3}$ | $\boldsymbol{v}$ | non－chronological | record trace | dynamic |
| ACD $^{4}$ | $\boldsymbol{v}$ | chronological | conjoin DSOPs | dynamic |

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## Contributions

- Abstract CNF2dDNNF - a calculus for CNF-to-d-DNNF compilation
- Combination of model enumeration with chronological CDCL with dynamic component analysis
- Formal correctness proof
- First CNF-to-d-DNNF compiler based on chronological CDCL


## Outline of the Rest of the Talk

The working of Abstract CNF2dDNNF

Abstract CNF2dDNNF by an example

Recap and ideas for further work

The Main Idea


The Main Idea


From DSOP to d-DNNF

$$
\begin{aligned}
F & =(\neg a \vee b \vee c) \wedge(\neg a \vee d) \wedge(a \vee b) \\
\operatorname{DSOP}(F) & =(\underline{a} \wedge b \wedge \underline{d}) \vee(\underline{a} \wedge \neg b \wedge c \wedge \underline{d}) \vee(\neg a \wedge b)
\end{aligned}
$$

From DSOP to d-DNNF

$$
\begin{gathered}
F=(\neg a \vee b \vee c) \wedge(\neg a \vee d) \wedge(a \vee b) \\
\operatorname{DSOP}(F)=(\underline{a} \wedge b \wedge \underline{d}) \vee(\underline{a} \wedge \neg b \wedge c \wedge \underline{d}) \vee(\neg a \wedge b)
\end{gathered}
$$

$$
\mathrm{d}-\operatorname{DNNF}(F)=(\underline{a} \wedge(b \vee(\neg b \wedge c)) \wedge \underline{d}) \vee(\neg a \wedge b)
$$

## CNF-to-d-DNNF Compilation

$$
(\neg a \vee b \vee c) \wedge(\neg a \vee d) \wedge(a \vee b) \quad 0
$$

## CNF－to－d－DNNF Compilation



## CNF-to-d-DNNF Compilation



## CNF-to-d-DNNF Compilation



## CNF-to-d-DNNF Compilation



## CNF-to-d-DNNF Compilation



## CNF-to-d-DNNF Compilation



## CNF-to-d-DNNF Compilation

$$
(a \wedge(b \vee(\neg b \wedge c)) \wedge d) \vee(\neg a \wedge b) \quad 0
$$

## Model Counting in d-DNNF

(1) $F=G \wedge H \quad \Longrightarrow \quad \# F=\# G \cdot \# H$
provided $\operatorname{var}(G) \cup \operatorname{var}(H)=\operatorname{var}(F)$ and $\operatorname{var}(G) \cap \operatorname{var}(H)=\emptyset$
(2) $F=C \vee D \Longrightarrow \# F=2^{|\operatorname{var}(F)|-|\operatorname{var}(C)|}+2^{|\operatorname{var}(F)|-|\operatorname{var}(D)|}$ provided $C \wedge D \equiv \perp$

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(1) $\quad F=(a \vee b) \wedge(c \vee d) \Longrightarrow \# F=3 \cdot 3=9$
$\mathcal{M}(F)=\{a b c d, a b c \neg d, a b \neg c d, a \neg b c d, a \neg b c \neg d, a \neg b \neg c d, \neg a b c d, \neg a b c \neg d, \neg a b \neg c d\}$

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$$

(2) $\quad F=a \vee(\neg a \wedge b \wedge c) \Longrightarrow \# F=2^{2}+2^{0}=5$
$\mathcal{M}(F)=\{a b c, a b \neg c, a \neg b c, a \neg b \neg c, \neg a b c\}$

Model Counting in d-DNNF

$$
\begin{align*}
F & =\underline{(a \wedge(b \vee(\neg b \wedge c)) \wedge d)} \vee \underline{(\neg a \wedge b)}  \tag{2}\\
\# F & =\# \underline{(a \wedge} \wedge \underline{(b \vee(\neg b \wedge c)}) \wedge \underline{d}) \cdot 2^{0}+\#(\underline{\neg a} \wedge \underline{b}) \cdot 2^{2}  \tag{1}\\
& =[\#(a) \cdot \# \underline{(b \vee} \underline{(\neg b \wedge c)}) \cdot \#(d)] \cdot 2^{0}+[\#(\neg a) \cdot \#(b)] \cdot 2^{2}  \tag{2}\\
& =\left[1 \cdot\left[\#(b) \cdot 2^{1}+\#(\underline{\neg b} \wedge \underline{c})\right] \cdot 2^{0}\right]+[1 \cdot 1] \cdot 2^{2}  \tag{1}\\
& =\left[1 \cdot\left[1 \cdot 2^{1}+[\#(\neg b) \cdot \#(c)] \cdot 2^{0}\right]+2^{2}\right. \\
& =\left[1 \cdot\left[1 \cdot 2^{1}+[1 \cdot 1] \cdot 2^{0}\right]+2^{2}=7\right.
\end{align*}
$$

## This work

## Ideas for future work

- CNF-to-d-DNNF compilation calculus
- Enumerative approach
- Based on chronological CDCL
- Formal proof of correctness
- Implementation (of proof of concept)
- Target compact formula representation
- Effect of decomposition on dual reasoning
- Investigate decomposability of formulae


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$$
\text { Thank you for your attention })
$$


[^0]:    ${ }^{1}$ Darwiche，ECAI＇04
    ${ }^{2}$ Muise et al．，CANAI＇12
    ${ }^{3}$ Lagniez \＆Marquis，IJCAl＇18
    ${ }^{4}$ Abstract CNF2dDNNF－this work

